

Clustering in Peer-to-Peer File Sharing Workloads

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Abstract

Peer-to-peer file sharing systems now generate a significant portion of Internet traffic. A good understanding of their workloads is crucial in order to improve their scalability, robustness and performance. Previous measurement studies on Kazaa and Gnutella were based on monitoring peer requests, and mostly concerned with peer and file availability and network traffic. In this paper, we take different measurements: instead of passively recording requests, we actively probe peers to get their cache contents information. This provides us with a map of contents, that we use to evaluate the degree of clustering in the system¹, and that could be exploited to improve significantly the search process.

1 Introduction

File sharing using peer-to-peer (P2P) networks has gained wide popularity; some reports [11, 13] suggest that P2P traffic is the dominant consumer of bandwidth ahead of Web traffic. This popularity triggered a lot of research activity. While one research trend aims at improving the performance and features [12, 17, 6], another trend is concerned with analysing and modeling these networks.

In the last few years, many measurement analysis of P2P networks focusing on *Free-Riding* [3], peer connectivity and availability (e.g., [4, 15, 14]), peer distribution and locality within the network (in [15, 9]) were carried out. But to the best of our knowledge no detailed analysis

¹We define clustering here as the overlap between cache contents.

has been done on the type of content shared on these networks.

Recent studies [9] show that locality-awareness in P2P networks can improve their efficiency significantly: by clustering users according to their *geographical* locality, file requests can be answered faster.

Similarly, clustering based on the types of files that peers have, also called *interest-based* locality, may improve search performance: it may reduce the duration of the search phase in general, and more specifically when searching for rare files [16]. Obviously, the performance gains obtained from exploiting either type of locality will depend on the corresponding degree of clustering. Our work investigates these issues.

Contribution. In this paper we present an analysis of contents (the type of files: music, video, documents etc) that peers offer to others in P2P file sharing networks. We used an active probing technique which allowed to capture the lists of files that peers offer. The community that we probed used eDonkey 2000, Kazaa's main competitor among all the P2P networks [10]. We then examine the clustering properties of the observed workload. Our preliminary findings show that both geographical and interest-based clustering are present and could thus be leveraged in order to yield significant performance improvements.

2 Experimental Setup

2.1 The eDonkey Network

The eDonkey 2000 network is one of the most advanced P2P file-sharing networks, providing features such as concurrent download of a file from

multiple sources, detection of file corruption using hashing, partial sharing of files while being downloaded and expressive querying methods for file search. According to a recent study [19], in some European countries eDonkey is now ahead of Kazaa in terms of resulting traffic.

The eDonkey network is a hybrid two-layer network composed of clients and servers. The clients connect to a server and register the files that they are sharing, providing the meta-data describing the files. Once registered, the clients can either search by querying meta-data or request a particular file via its unique network identifier. Servers provide the locations of files when requested by clients, so that clients can download the files directly from the provided locations.

2.2 Methodology

Measurements reported in this paper have been obtained by crawling the eDonkey network during the first week of November 2003. 230,000 eDonkey clients were discovered; 55,000 of them were connected during a 3 days period. We kept 37,000 clients that could clearly be identified as distinct clients, among which 25,000 clients shared no files at all. Our study is based on the 923,000 different files shared by the remaining 12,000 clients.

Our crawler has been implemented by modifying an existing eDonkey client, namely MLdonkey [8]. Our modified client runs two concurrent tasks: discovering eDonkey clients and scanning their contents.

Client Discovery. Our crawler first connects to as many eDonkey servers as possible requesting their list of clients (server responses are limited to 200 clients). In order to obtain as many clients as possible from each server (servers may have up to 100,000 simultaneously connected users), we send requests for clients containing the strings “aa”, “ab”, “ac”, ... and so on until “zz”.

Client Content Scanning. Our crawler attempts to connect to every eDonkey client that is discovered. If it succeeds, it obtains the unique identifier of the client and requests its list of shared files. All these data are then stored in

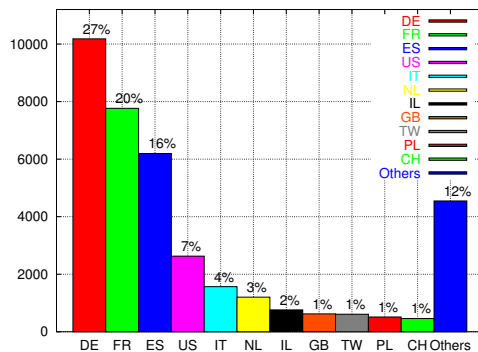


Figure 1: Distribution of peers per country

a file, named using a digest of the IP address and port of the client, so that the resulting trace is completely anonymised.

3 Workload Properties

Our crawler provides us with the content lists of each peer that replied with a non empty list of files. The first set of results presents the properties of the workload in terms of popularity, distribution between various types of files and sharing profiles of peers.

As opposed to previous studies which gathered requests from a particular location – such as a university [9]–, we gather a map of contents over a large number of countries (see Figure 1).

3.1 Replication Patterns

While previous studies evaluated the popularity as the number of requests per file, we measure a file’s popularity by its replication degree. Results match former observations on file popularity; *i.e.*, a few files are extremely replicated while a large number is not replicated at all and approximately 100,000 files are present in two caches or more. Figures 2 and 3 respectively present the replica distribution of files for all files and for various types of files, per file rank. We observe similar properties between the eDonkey workloads and the Kazaa workloads presented in [9], and in particular after an initial flat region, the popularity distributions have a clear linear trend on log-log plots. Therefore, whether the popularity of a file is expressed in terms of

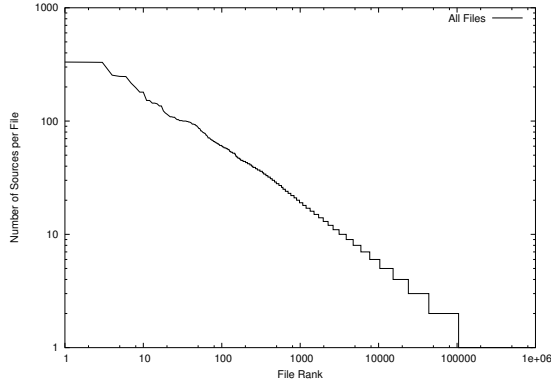


Figure 2: Replica distribution of files (900,000 files)

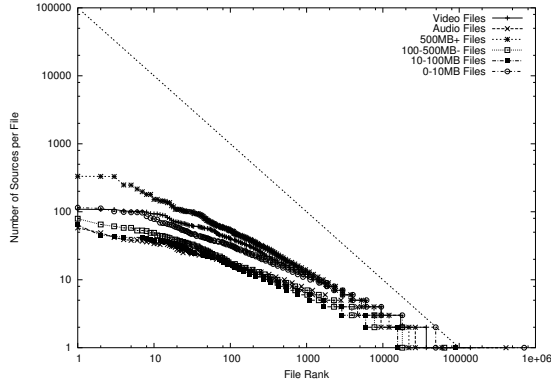


Figure 3: Replica distribution of files depending of their type

number of requests or number of replicas, the observed patterns are similar.

3.2 Sharing and Distribution

Figures 4 and 5 display respectively the CDF of number of files and the disk-space that each peer offers to the network. While the proportion of free-riding is still significant (68 %), most of the remaining clients share few files (between 10 and 100), but large files (between 1 and 50 GB of shared data, the protocol limit for file size is 4 GB).

We have also investigated the proportion of files of a given type share in the system, measured both in number of files and disk-space. We observe that multimedia files, audio and video, dominate: Audio files represent the largest number of files (48% against 16% for video), while in terms of size video files are dominant (67%

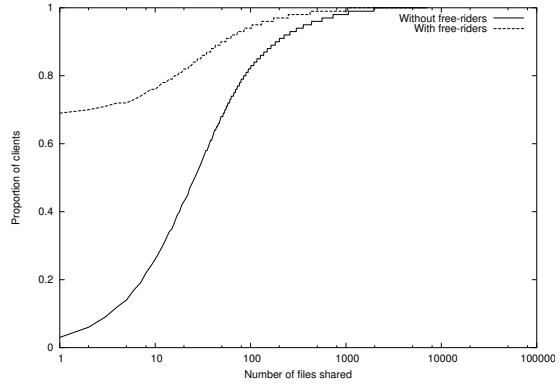


Figure 4: Distribution of the number of files *offered* per client

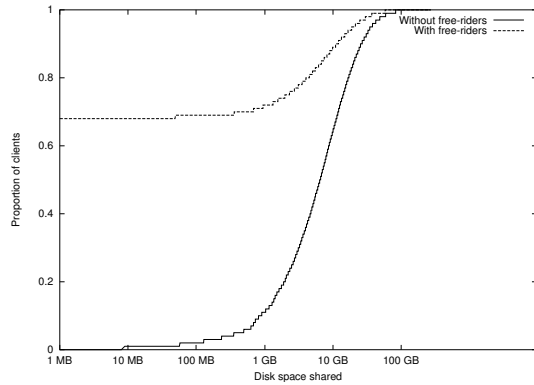


Figure 5: Distribution of the disk-space shared per client

against 16% for audio). Other measures show that *MP3* files dominate in number, while *AVI* files largely dominate in terms of size. Although such distribution is expected and matches results obtained on different workloads, it is interesting since it leads to the exploitation of two different types of clustering: geographical and interest-based clustering.

3.3 Clustering

For clarity, we consider only the two major types of files, audio and video files, leading to different opportunities in terms of clustering exploitation.

Geographical Clustering and Video Files.

It seems reasonable to assume that the traffic generated by video files dominates largely in the network. For such files, the latency and traf-

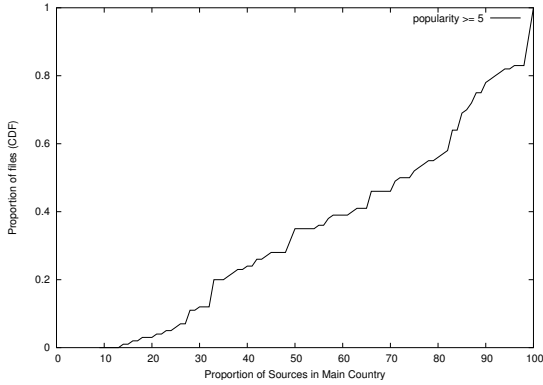


Figure 6: Proportion of sources in one country (for video files with more than 5 sources).

fic induced by the search process are negligible as compared to the download phase. Figure 6 shows a measure of geographical clustering for video files in the collected workload. The graph can be read as follows: the y -value of a point represents the proportion of video files (having more than 5 replicas) for which x -value% of the sources are located in the main country of replication². For example, the graph shows that, for 60 % of the files, more than 80% of the replicas are located in the main country. This definitively attests of the presence of geographical clustering, *i.e.*, peers requesting a given video file may in a large proportion of cases download it from peers in their own country, thus achieving low latency and network usage compared to downloading it from a randomly chosen peer.

Interest-based Clustering and Audio Files.

As 48 % of the shared files are audio files, it is reasonable to assume that most requests are for such files.

Performance of P2P search for such files can be greatly improved if we exploit interest-based locality [16, 7, 18]: if two peers share interests, (in other words if the contents of their cache overlap significantly) the search mechanism can be significantly improved by having these peers connect to one another and first send their requests to each other.

²The main country is defined as the country where the majority of replicas are hosted.

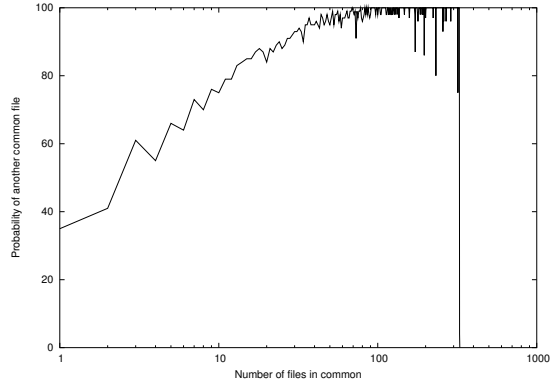


Figure 7: Clustering between peers (for all files)

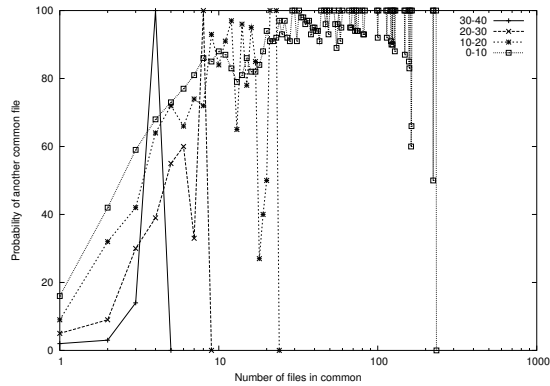


Figure 8: Clustering between peers for audio files (according to popularity)

Clustering Measurements. Figure 7 displays the clustering between every pair of peers and is measured as the probability that any two clients having at least a given number of files in common share (at least) another one. We observe that this curve is increasing very quickly with the number of files in common, up to a certain point (here, 325) that reflects the maximal cache overlap in the observed trace. As soon as some peers have a small number of files in common (say 10), the probability is high (approximately 0.8) that they will have another one in common. This probability is very close to 1 for peers sharing over 50 files in common.

Figure 8 presents the same clustering information, calculated for sub-classes of files, namely for audio files, and depending on their popularity. Results indicate that the clustering is pretty high for un-popular audio files (2-10 replicas) and in addition the maximal overlap between

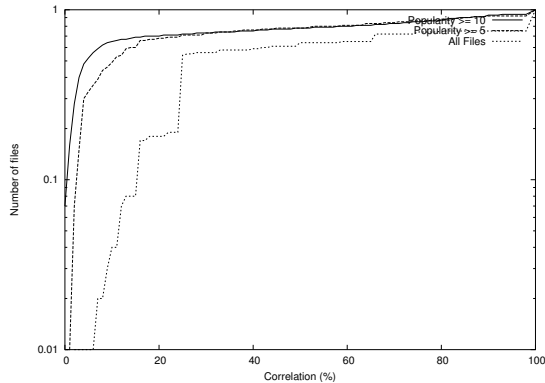


Figure 9: The CDF of maximal correlation to other files (for distinct file popularities)

peers cache contents reach over 200 unpopular audio files. As the popularity increases, the overlap tends to decrease as well (30-40 replicas). However, the probability increases consistently and fairly quickly with the number of files in common.

We also evaluated the correlation between files, rather than between peers. We define the correlation between two files as the correlation coefficient between the indicator variables equal to 1 if a randomly chosen peer has the document, and zero otherwise. Figure 9 displays the cumulative distribution function of the maximal correlation coefficients of files to other files. For example a correlation value of 60 % on the x-axis corresponds to a y-value of less than 70 % (for the curve of all files). Thus, at least 30 % of the files have a correlation larger than 60 % to some other file. This can be expected, as many files are cdrom images of programs or movies in several parts, sets of cracks for programs, and un-popular MP3 tracks of popular albums.

Further interpretation of the results, as well as processing of the data is needed to get a better characterization of interest-based clustering and how it might be exploited. Such extensions are currently under way.

4 Conclusion

Peer to peer file sharing is now the main component of Internet traffic. Yet, there have been

only a few studies of workload measurements and even less consideration of clustering properties. In this paper, we presented the results of an evaluation we conducted on the most popular file sharing network in Europe, namely the eDonkey 2000 network, in an attempt to provide a map of the content shared by its peers. In particular, we focused on the geographical and the interest-based clustering properties for video and audio files respectively, of the observed workload. We demonstrated that there is a significant locality of interest for audio files, that could be used to improve the search mechanisms without relying on servers. Interest-based locality can be exploited to create additional links to semantically related peers. These peers are contacted first before sending a request to the servers or to flood the request over the network. The advantage of implementing these additional links is that the technique is generic and can be applied to both flooding-based peer-to-peer networks, whether they are based on a structured [5] or unstructured network [1], and server-based [2, 8].

The interest of this measurement study is three fold: First, to the best of our knowledge, it represents the first evaluation study focusing on clustering properties and providing such a map of contents. Second, we also aim at doing a more comprehensive study by looking at a broader spectrum of users from different countries, and not limited to a particular community. Finally, the workloads we evaluated can be used as input to evaluate peer-to-peer file sharing systems and contribute to the improvement of such systems.

The trace collected demonstrates a certain degree of clustering between peers, we are currently investigating the impact of the generous peers (both in space and number of files) to evaluate their impact on the detected semantic relationships.

Analysing workloads of peer-to-peer file sharing system is crucial to understand and improve such systems and is very challenging as it depends on the cooperation of peers. Future work includes a similar study for the Kazaa network in order to study the similarities and differences between the two networks.

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