Integrating supply and demand within the framework of mixed integer linear problems

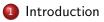
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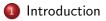
# Outline











2) General framework

3 Case study



# Motivation

#### Demand

- Choices of customers
- Discrete choice models
- Nonlinear and nonconvex formulations

### Supply

- Design and configuration of the system
- Mixed Integer Linear Problems (MILP)

### Demand model



- Population of N customers (n)
- Choice set C(i)
- $C_n \subseteq C$ : alternatives considered by customer n $(\mathcal{N}_i = \{n \ge 1 | i \in C_n\})$

Behavioral assumption

• 
$$U_{in} = V_{in} + \varepsilon_{in}$$

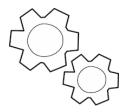
• 
$$V_{in} = \sum_{k} \beta_{ink} x^{e}_{ink} + q^{d}(x^{d})$$
  
•  $P_{n}(i|\mathcal{C}_{n}) = \Pr(U_{in} \ge U_{jn}, \forall j \in \mathcal{C}_{n})$ 

### Simulation

- Distribution  $\varepsilon_{in}$
- R draws  $\xi_{in1}, \ldots, \xi_{inR}$

• 
$$U_{inr} = V_{in} + \xi_{inr}$$

# Supply model



- Operator selling services to a market
  - Price *p*<sub>in</sub> (to be decided)
  - Capacity c<sub>i</sub>
- Benefit (revenue cost) to be maximized
- Opt-out option (*i* = 0)

### Price characterization

- Lower and upper bound
- Discretization: price levels
- Binary representation  $(\lambda_{in\ell})$

#### Capacity allocation

- Exogenous priority list of customers
- Here it is assumed as given
- Capacity as decision variable



### 2 General framework

3 Case study



# MILP (in words)

### MILP

max benefit subject to availability utility definition discounted utility choice capacity allocation price selection

### MILP

max benefit subject to **availability** utility definition discounted utility choice capacity allocation price selection  $\begin{array}{ll} y_i \in \{0,1\} & \text{operator decision} \\ y_{in}^d \in \{0,1\} & \text{customer decision (data)} \\ y_{in} \in \{0,1\} & \text{product of decisions} \\ y_{inr} \in \{0,1\} & \text{capacity restrictions} \end{array}$ 

#### Relations between availabilities

$$y_{in} = y_{in}^{d} y_{i} \quad \forall i, n$$
(1)  
$$y_{inr} \leq y_{in} \quad \forall i, n, r$$
(2)

#### General framework

# MILP

### MILP

max benefit subject to availability utility definition discounted utility choice capacity allocation price selection

$$U_{inr} \qquad \text{utility}$$

$$z_{inr} = \begin{cases} U_{inr} & \text{if } y_{inr} = 1\\ \ell_{nr} & \text{if } y_{inr} = 0 \end{cases} \quad \text{discounted utility}$$

$$(\ell_{nr} \text{ smallest lower bound})$$
Utility
$$U_{inr} = \overbrace{\beta_{in}p_{in} + q_d(x_d)}^{V_{in}} + \xi_{inr} \forall i, n, r \qquad (3)$$

#### Discounted utility

$$\ell_{nr} \leq z_{inr} \qquad \forall i, n, r \quad (4)$$

$$z_{inr} \leq \ell_{nr} + M_{inr}y_{inr} \quad \forall i, n, r \quad (5)$$

$$U_{inr} - M_{inr}(1 - y_{inr}) \le z_{inr} \qquad \forall i, n, r \quad (6)$$

$$z_{inr} \leq U_{inr}$$
  $\forall i, n, r$  (7)

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### MILP

max benefit subject to availability utility definition discounted utility **choice** capacity allocation price selection

$$U_{nr} = \max_{i \in C} z_{inr}$$
$$w_{inr} = \begin{cases} 1 & \text{if } i = \arg \max\{U_{nr}\}\\ 0 & \text{otherwise} \end{cases}$$
 choice

#### Choice

$$z_{inr} \leq U_{nr}$$
  $\forall i, n, r$  (8)

$$U_{nr} \leq z_{inr} + M_{nr}(1 - w_{inr}) \qquad \forall i, n, r \qquad (9)$$

$$\sum_{i} w_{inr} = 1 \qquad \qquad \forall n, r \qquad (10)$$

$$w_{inr} \leq y_{inr}$$
  $\forall i, n, r$  (11)

#### MILP

max benefit subject to availability utility definition discounted utility choice **capacity allocation** price selection

#### Priority list

$$y_{in^{-}r} \ge y_{inr} \qquad \forall i > 0, n < N, r \qquad (12)$$

Capacity cannot be exceeded  $\Rightarrow y_{inr} = 1$ 

$$\sum_{m=1}^{n-1} w_{imr} \le (c_i - 1) y_{inr} + (n-1)(1 - y_{inr}) \quad \forall i > 0, n > c_i, r \quad (13)$$

Capacity has been reached  $\Rightarrow y_{inr} = 0$ 

$$c_i(y_{in}-y_{inr}) \leq \sum_{m=1}^{n-1} w_{imr} \quad \forall i > 0, n, r \quad (14)$$

#### MILP

max benefit subject to availability utility definition discounted utility choice capacity allocation **price selection** 

$$p_{in} = \frac{1}{10^k} \left( \ell_{in} + \sum_{\ell=0}^{L_{in}-1} 2^\ell \lambda_{in\ell} \right)$$

When calculating the benefit: λ<sub>inℓ</sub>w<sub>inr</sub>
 α<sub>inℓ</sub> = λ<sub>inℓ</sub>w<sub>inr</sub>

Linearization of  $\alpha_{inr\ell}$  + Price bounded from above

$$\lambda_{in\ell} + w_{inr} \leq 1 + \alpha_{inr\ell} \quad \forall i > 0, n, r, \ell \quad (15)$$
  
$$\alpha_{inr\ell} \leq \lambda_{in\ell} \qquad \forall i > 0, n, r, \ell \quad (16)$$
  
$$\alpha_{inr\ell} \leq w_{inr} \qquad \forall i > 0, n, r, \ell \quad (17)$$

$$\ell_{in} + \sum_{\ell=0}^{L_{in}-1} 2^{\ell} \lambda_{in\ell} \leq m_{in} \qquad \forall i > 0, n \qquad (18)$$

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### MILP

max benefit subject to availability utility definition discounted utility choice capacity allocation price selection

$$\max \sum_{i>0} (R_i - C_i)$$

#### Revenue

$$R_{i} = \frac{1}{R} \frac{1}{10^{k}} \left[ \sum_{n} \sum_{r} \left( \ell_{in} w_{inr} + \sum_{\ell} 2^{\ell} \alpha_{inr\ell} \right) \right]$$

### Cost

$$C_i = (f_i + v_i c_i) y_i$$

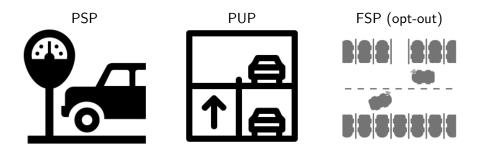


### 2 General framework





# Parking choices<sup>1</sup>



- N = 50 customers
- $C = \{PSP, PUP, FSP\}$
- $C_n = C \quad \forall n$

- $p_{in} = p_i \quad \forall n$
- Mixtures of a logit model

<sup>1</sup>A. Ibeas, L. dellOlio, M. Bordagaray, <u>et al.</u>, "Modelling parking choices considering user heterogeneity," <u>Transportation Research Part A: Policy and Practice</u>, vol. 70, pp. 41 –49, 2014. MP, SSA, MB, BG IFORS 2017 11 / 18

# General experiments

#### Uncapacitated vs Capacitated case

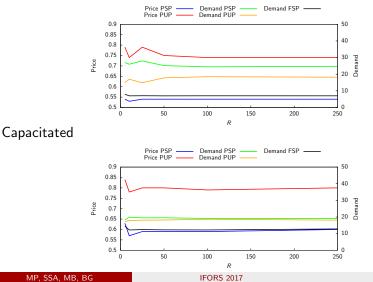
- Maximization of revenue
- Unlimited capacity
- Capacity of 20 spots for PSP and PUP

#### Price differentiation by population segmentation

- Reduced price for residents
- Two scenarios
  - Subsidy offered by the municipality
  - Operator is obliged to offer a reduced price

# Uncapacitated vs Capacitated case

Uncapacitated



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# Computational time

|     | Uncapacitated case |      |      |       | Capacitated case |      |      |       |
|-----|--------------------|------|------|-------|------------------|------|------|-------|
| R   | Sol time           | PSP  | PUP  | Rev   | Sol time         | PSP  | PUP  | Rev   |
| 5   | 2.58 s             | 0.54 | 0.79 | 26.43 | 12.0 s           | 0.63 | 0.84 | 25.91 |
| 10  | 3.98 s             | 0.53 | 0.74 | 26.36 | 54.5 s           | 0.57 | 0.78 | 25.31 |
| 25  | 29.2 s             | 0.54 | 0.79 | 26.90 | 13.8 min         | 0.59 | 0.80 | 25.96 |
| 50  | 4.08 min           | 0.54 | 0.75 | 26.97 | 50.2 min         | 0.59 | 0.80 | 26.10 |
| 100 | 20.7 min           | 0.54 | 0.74 | 26.90 | 6.60 h           | 0.59 | 0.79 | 26.03 |
| 250 | 2.51 h             | 0.54 | 0.74 | 26.85 | 1.74 days        | 0.60 | 0.80 | 25.93 |

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### 2 General framework

3 Case study



# Lagrangian relaxation

#### General idea

- Decompose the MILP into 2 subproblems
- Solve the subproblems independently
- Lagrangian dual to provide an upper bound

#### Operator subproblem

• Resulting problem: Capacitated Facility Location Problem

#### Customer supbroblem

- Assumption: utility decreases as a function of the price
- Iterate over customers (priority list) and over scenarios
- Highest price such that the customer does not change the choice

# Ongoing research and future work

### Ongoing research

- Implementation of the 2 subproblems
- Subgradient method to solve the Lagrangian dual

#### Future work

- Provide a lower bound on the original problem
- If the gap between bounds is significant  $\Rightarrow$  column generation

### Questions?



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