

# The automorphism group of a self-dual [72, 36, 16] code is not an elementary abelian group of order 8

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## Abstract

The existence of an extremal self-dual binary linear code  $\mathcal{C}$  of length 72 is a long-standing open problem. We continue the investigation of its automorphism group: looking at the combination of the subcodes fixed by different involutions and doing a computer calculation with MAGMA, we prove that  $\text{Aut}(\mathcal{C})$  is not isomorphic to the elementary abelian group of order 8. Combining this with the known results in the literature one obtains that  $\text{Aut}(\mathcal{C})$  has order at most 5.

*Keywords:* automorphism group, self-dual extremal codes

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## 1. Introduction

A binary linear code of length  $n$  is a subspace of  $\mathbb{F}_2^n$ , where  $\mathbb{F}_2$  is the field with 2 elements. A binary linear code  $\mathcal{C}$  is called *self-dual* if  $\mathcal{C} = \mathcal{C}^\perp$  with respect to the Euclidean inner product. It follows immediately that the dimension of such a code has to be the half of the length. The *minimum distance* of  $\mathcal{C}$  is defined as  $d(\mathcal{C}) := \min_{c \in \mathcal{C} \setminus \{0\}} \#\{i \mid c_i = 1\}$ . In [7] an upper bound for the minimum distance of self-dual binary linear codes is given. Codes achieving this bound are called *extremal*. The most interesting codes, for various reasons, are those whose length is a multiple of 24: in this case  $d(\mathcal{C}) = 4m + 4$ , where  $24m$  is the length of the code, and they give rise to beautiful combinatorial structures [2]. There are unique extremal self-dual codes of length 24 (the extended binary Golay code  $\mathcal{G}_{24}$ ) and 48 (the extended quadratic residue code  $\mathcal{QR}_{48}$ ). For nearly forty years many people have tried

unsuccessfully to find an extremal self-dual code of length 72 [9]. The usual approach to this problem is to study the possible automorphism groups (see next section for the detailed definition of it). Most of the subgroups of  $S_{72}$  are now excluded: the last result is contained in [4], in which the authors finished to exclude all the non-abelian groups with order greater than 5.

In this paper we prove that the elementary abelian group of order 8 cannot occur as automorphism group of such a code, obtaining the following.

**Theorem 1.1.** *The automorphism group of a self-dual [72, 36, 16] code is either cyclic of order 1, 2, 3, 4, 5 or elementary abelian of order 4.*

The techniques which we use are similar to those of [3]. We know [8], up to equivalence, the possible subcodes fixed by all the non-trivial involutions. So we combine them pairwise, checking the minimum distance to be 16, and we classify their sum, up to equivalence. We get only a few extremal codes and all of them satisfy certain intersection properties that, with easy dimension arguments, make it impossible to sum a third fixed subcode without loosing the extremality.

All results are obtained using extensive computations in MAGMA [5].

## 2. Basic definitions and notations

Throughout the paper we will use the following notations for groups:

- $C_m$  is the cyclic group of order  $m$ ;
- $S_m$  is the symmetric group of degree  $m$ ;
- if  $A$  and  $B$  are two groups,  $A \times B$  indicates their direct product;
- if  $A$  and  $B$  are two groups,  $A \wr B$  indicates their wreath product.

Given a group  $G$  and a subgroup  $H$  of  $G$  ( $H \leq G$ ) we denote  $C_G(H)$  the centralizer of  $H$  in  $G$ . Let  $\kappa \in G$ . Then  $C_G(\kappa) := C_G(\langle \kappa \rangle)$ , where  $\langle \kappa \rangle$  is the (cyclic) group generated by  $\kappa$ .

Let us consider the ambient space  $\mathbb{F}_2^n$ . We will indicate with calligraphic capital letters the subspaces of  $\mathbb{F}_2^n$ , in order to distinguish them from groups. We have a natural (right) action of  $S_n$  on  $\mathbb{F}_2^n$  defined as follows: let  $v = (v_1, \dots, v_n) \in \mathbb{F}_2^n$  and  $\sigma \in S_n$ ; then

$$v^\sigma := (v_{1\sigma^{-1}}, \dots, v_{n\sigma^{-1}}).$$

We have an action induced naturally on the subspaces of  $\mathbb{F}_2^n$ :

$$\mathcal{C}^\sigma := \{c^\sigma \mid c \in \mathcal{C}\},$$

where  $\mathcal{C} \leq \mathbb{F}_2^n$  and  $\sigma \in S_n$ .

Let  $\mathcal{C} \leq \mathbb{F}_2^n$ . Then the *automorphism group* of the code  $\mathcal{C}$  is the subgroup of  $S_n$  defined as

$$\text{Aut}(\mathcal{C}) := \{\sigma \in S_n \mid \mathcal{C}^\sigma = \mathcal{C}\}.$$

Given a code  $\mathcal{C}$  and an automorphism  $\sigma \in \text{Aut}(\mathcal{C})$  we define

$$\mathcal{C}(\sigma) := \{c \in \mathcal{C} \mid c^\sigma = c\}.$$

This is a subcode of  $\mathcal{C}$  and we call it the *subcode fixed by  $\sigma$* .

### 3. Preliminary observations

Let  $\mathcal{C}$  be a self-dual [72, 36, 16] code such that

$$\text{Aut}(\mathcal{C}) \cong C_2 \times C_2 \times C_2 = \langle \alpha, \beta, \gamma \rangle.$$

By [6] all non-trivial elements of  $\text{Aut}(\mathcal{C})$  are fixed point free (that is of degree  $n$ ) and we may relabel the coordinates so that

$$\begin{aligned} \alpha &= (1, 2)(3, 4)(5, 6)(7, 8) \dots (71, 72) \\ \beta &= (1, 3)(2, 4)(5, 7)(6, 8) \dots (70, 72) \\ \gamma &= (1, 5)(2, 6)(3, 7)(4, 8) \dots (68, 72). \end{aligned}$$

**Definition 3.1.** Let  $\mathcal{V} := \mathbb{F}_2^n$ . Then

$$\begin{aligned} \pi_\alpha : \mathcal{V}(\alpha) &\rightarrow \mathbb{F}_2^{36} \\ (v_1, v_1, v_2, v_2, v_3, v_3, v_4, v_4, \dots, v_{36}, v_{36}) &\mapsto (v_1, v_2, \dots, v_{36}) \end{aligned}$$

denote the bijection between the subspace of fixed by  $\alpha$  and  $\mathbb{F}_2^{36}$ ,

$$\begin{aligned} \pi_\beta : \mathcal{V}(\beta) &\rightarrow \mathbb{F}_2^{36} \\ (v_1, v_2, v_1, v_2, v_3, v_4, v_3, v_4, \dots, v_{35}, v_{36}) &\mapsto (v_1, v_2, \dots, v_{36}) \end{aligned}$$

denote the bijection between the subspace fixed by  $\beta$  and  $\mathbb{F}_2^{36}$  and

$$\begin{aligned} \pi_\gamma : \mathcal{V}(\gamma) &\rightarrow \mathbb{F}_2^{36} \\ (v_1, v_2, v_3, v_4, v_1, v_2, v_3, v_4, \dots, v_{35}, v_{36}) &\mapsto (v_1, v_2, \dots, v_{36}) \end{aligned}$$

denote the bijection between the subspace fixed by  $\gamma$  and  $\mathbb{F}_2^{36}$ .

**Remark 3.2.** The centralizer  $C_{S_{72}}(\alpha) \cong C_2 \wr S_{36}$  of  $\alpha$  acts on the set of fixed points of  $\alpha$ . Using the isomorphism  $\pi_\alpha$  we hence obtain a group epimorphism which we denote by  $\eta_\alpha$

$$\eta_\alpha : C_{S_{72}}(\alpha) \rightarrow S_{36}$$

with kernel  $C_2^{36}$ . Similarly we obtain the epimorphisms

$$\eta_\beta : C_{S_{72}}(\beta) \rightarrow S_{36}$$

and

$$\eta_\gamma : C_{S_{72}}(\gamma) \rightarrow S_{36}.$$

By [8] we have that all the projections of the fixed codes  $\pi_\alpha(\mathcal{C}(\alpha)), \pi_\beta(\mathcal{C}(\beta))$  and  $\pi_\gamma(\mathcal{C}(\gamma))$  are self-dual [36, 18, 8] codes. Such codes have been classified in [1], up to equivalence (under the action of the full symmetric group  $S_{36}$ ) there are 41 such codes. Notice that

$$\langle \eta_\alpha(\beta), \eta_\alpha(\gamma) \rangle = \langle \eta_\beta(\alpha), \eta_\beta(\gamma) \rangle = \langle \pi_\gamma(\alpha), \eta_\gamma(\beta) \rangle = \langle \chi, \mu \rangle \leq S_{36},$$

with

$$\chi = (1, 2)(3, 4) \dots (35, 36)$$

and

$$\mu = (1, 3)(2, 4) \dots (34, 36),$$

are contained in  $\text{Aut}(\pi_\alpha(\mathcal{C}(\alpha))), \text{Aut}(\pi_\beta(\mathcal{C}(\beta)))$  and  $\text{Aut}(\pi_\gamma(\mathcal{C}(\gamma)))$  respectively. Only 14 of the 41 codes, say  $\mathbb{Y} := \{\mathcal{Y}_1, \dots, \mathcal{Y}_{14}\}$ , have an automorphism group which contains at least one subgroup conjugate to  $\langle \chi, \mu \rangle$ .

By direct calculation on these 14 codes we get the following conditions on the intersection of the codes.

**Lemma 3.3.** Let

$$(\chi', \mu', \zeta') \in \{(\alpha, \beta, \gamma), (\alpha, \gamma, \beta), (\beta, \alpha, \gamma), (\beta, \gamma, \alpha), (\gamma, \alpha, \beta), (\gamma, \beta, \alpha)\}.$$

Then we have only the following possibilities:

| $\dim(\mathcal{C}(\chi') \cap \mathcal{C}(\mu') \cap \mathcal{C}(\zeta'))$ | $\dim(\mathcal{C}(\chi') \cap \mathcal{C}(\mu'))$ | $\dim(\mathcal{C}(\chi') \cap \mathcal{C}(\zeta'))$ |
|--|---|---|
| 5  | 9   | 9   |
| 5  | 9   | 10  |
| 6  | 9   | 9   |
| 6  | 9   | 10  |
| 6  | 9   | 11  |
| 6  | 10  | 10  |
| 6  | 10  | 11  |

Let  $G := C_{S_{72}}(\text{Aut}(\mathcal{C}))$ . Then  $G$  acts on the set of extremal self-dual codes with automorphism group  $\langle \alpha, \beta, \gamma \rangle$  and we aim to find a system of orbit representatives for this action. Here we have some differences with the non-abelian cases, since the full group  $\langle \alpha, \beta, \gamma \rangle$  is a subgroup of the automorphism group of all the fixed subcodes  $\mathcal{C}(\alpha)$ ,  $\mathcal{C}(\beta)$  and  $\mathcal{C}(\gamma)$ . The main property that we use is the following, which is straightforward to prove:

$$\pi_\alpha(\mathcal{C}(\alpha))(\chi) = \pi_\beta(\mathcal{C}(\beta))(\chi) \quad (1)$$

and similar relations for the other fixed subcodes. This allows us to combine properly  $\mathcal{C}(\alpha)$  and  $\mathcal{C}(\beta)$  classifying their sum.

#### 4. Description of the calculations

Let

$$\mathbb{D} := \{\mathcal{D} = \mathcal{D}^\perp \leq \mathbb{F}_2^{36} \mid d(\mathcal{D}) = 8, \langle \chi, \mu \rangle \leq \text{Aut}(\mathcal{D})\}.$$

The group

$$G_{36} := C_{S_{36}}(\langle \chi, \mu \rangle) = \eta_\alpha(G) = \eta_\beta(G) = \eta_\gamma(G)$$

acts, naturally, on this set.

**Lemma 4.1.** *A set of representatives of the  $G_{36}$ -orbits on  $\mathbb{D}$  can be computed by performing the following computations on each  $\mathcal{Y} \in \mathbb{Y}$ :*

- Let  $\chi_1, \dots, \chi_s$  represent the conjugacy classes of fixed point free elements of order 2 in  $\text{Aut}(\mathcal{Y})$ .
- Compute elements  $\tau_1, \dots, \tau_s \in S_{36}$  such that  $\tau_k^{-1}\chi_k\tau_k = \chi$  and put  $\mathcal{Y}_k := \mathcal{Y}^{\tau_k}$  so that  $\chi \in \text{Aut}(\mathcal{Y}_k)$ .
- For every  $\mathcal{Y}_k$ , consider the set of fixed point free elements  $\tilde{\mu}$  of order 2 in  $C_{\text{Aut}(\mathcal{Y}_k)}(\chi)$  such that  $\langle \chi, \tilde{\mu} \rangle$  is conjugate to  $\langle \chi, \mu \rangle$  in  $S_{36}$ . Let  $\mu_1, \dots, \mu_{t_k}$  represent the  $C_{\text{Aut}(\mathcal{Y}_k)}(\chi)$ -conjugacy classes in this set.
- Compute elements  $\sigma_1, \dots, \sigma_{t_k} \in C_{S_{36}}(\chi)$  such that  $\sigma_l^{-1}\mu_l\sigma_l = \mu$  and put  $\mathcal{Y}_{k,l} := \mathcal{Y}_k^{\sigma_l}$  so that  $\langle \chi, \mu \rangle \leq \text{Aut}(\mathcal{Y}_{k,l})$ .

Then  $\mathbb{D}' := \{\mathcal{Y}_{k,l} \mid \mathcal{Y} \in \mathbb{Y}, 1 \leq k \leq s, 1 \leq l \leq t_k\}$  represents the  $G_{36}$ -orbits on  $\mathbb{D}$ .

*Proof.* Clearly these codes lie in  $\mathbb{D}$ .

Since  $G_{36} \leq S_{36}$ , if we consider different elements in  $\mathbb{Y}$ , say  $\mathcal{Y}$  and  $\mathcal{Y}'$ , then  $\mathcal{Y}'_{k',l'}$  is not in the same orbit of  $\mathcal{Y}_{k,l}$  for any  $k', l', k, l$ .

Now assume that there is some  $\lambda \in G_{36}$  such that

$$\mathcal{Y}^{\tau_{k'}\sigma_{l'}} = \mathcal{Y}_{k',l'}^\lambda = \mathcal{Y}_{k,l} = \mathcal{Y}^{\tau_k\sigma_l}.$$

Then

$$\epsilon := \tau_{k'}\sigma_{l'}\lambda\sigma_l^{-1}\tau_k^{-1} \in \text{Aut}(\mathcal{Y})$$

satisfies  $\epsilon\chi_k\epsilon^{-1} = \chi_{k'}$ , so  $\chi_k$  and  $\chi_{k'}$  are conjugate in  $\text{Aut}(\mathcal{Y})$ , which implies  $k = k'$  (and so  $\tau_k = \tau_{k'}$ ). Now,

$$\mathcal{Y}^{\tau_k\sigma_{l'}\lambda} = \mathcal{Y}_k^{\sigma_{l'}\lambda} = \mathcal{Y}_k^{\sigma_l} = \mathcal{Y}^{\tau_k\sigma_l}.$$

Then

$$\epsilon' := \sigma_{l'}\lambda\sigma_l^{-1} \in \text{Aut}(\mathcal{Y}_k)$$

commutes with  $\chi$ . Furthermore  $\epsilon'\sigma_l\epsilon'^{-1} = \sigma_{l'}$  and hence  $l = l'$ .

Now let  $\mathcal{Z} \in \mathbb{D}$  and choose some  $\xi \in S_{36}$  such that  $\mathcal{Z}^\xi = \mathcal{Y} \in \mathbb{Y}$ . Then  $\xi^{-1}\chi\xi$  is conjugate to some of the chosen representatives  $\chi_i \in \text{Aut}(\mathcal{Y})$  ( $i = 1, \dots, s$ ) and we may multiply  $\xi$  by some automorphism of  $\mathcal{Y}$  so that

$$\xi^{-1}\chi\xi = \chi_k = \tau_k\chi\tau_k^{-1}.$$

So  $\xi\tau_k \in C_{S_{36}}(\chi)$  and  $\mathcal{Z}^{\xi\tau_k} = \mathcal{Y}^{\tau_k} = \mathcal{Y}_k$ .

It is straightforward to prove that the element  $(\xi\tau_k)^{-1}\mu(\xi\tau_k) \in \text{Aut}(\mathcal{Y}_k)$  is a fixed point free element of order 2 in  $C_{\text{Aut}(\mathcal{Y}_k)}(\chi)$  such that  $\langle \chi, (\xi\tau_k)^{-1}\mu(\xi\tau_k) \rangle$  is conjugate to  $\langle \chi, \mu \rangle$  in  $S_{36}$ . So there is some automorphism  $\omega \in C_{\text{Aut}(\mathcal{Y}_k)}(\chi)$  and some  $l \in \{1, \dots, t_k\}$  such that  $(\xi\tau_k\omega)^{-1}\mu(\xi\tau_k\omega) = \mu_l$ . Then

$$\mathcal{Y}^{\xi\tau_k\omega\sigma_l} = \mathcal{Y}_{k,l}$$

where  $\xi\tau_k\omega\sigma_l \in G_{36}$ . □

There are 242 such representatives. For our purposes we need to modify this set a little: consider the set  $\{\mathcal{Y}(\chi) \mid \mathcal{Y} \in \mathbb{D}\}$  and take a set of representatives for the action of  $G_{36}$  on this set, say  $\mathbb{E} := \{\mathcal{E}_1, \dots, \mathcal{E}_m\}$ . By calculations  $m = 40$ . For every  $1 \leq i \leq m$  define the set

$$\tilde{\mathbb{D}}_i := \{\mathcal{Y}^\epsilon \mid \mathcal{Y} \in \mathbb{D}' \text{ such that there exists } \epsilon \in G_{36} \text{ so that } \mathcal{Y}(\chi)^\epsilon = \mathcal{E}_i\}.$$

Clearly  $\bigcup_{i=1}^m \tilde{\mathbb{D}}_i$  is still a set of representatives of the  $G_{36}$ -orbits on  $\mathbb{D}$ , but now  $\mathcal{Y}_j(\chi)$  and  $\mathcal{Y}_k(\chi)$  are equal if  $\mathcal{Y}_j$  and  $\mathcal{Y}_k$  belong to the same  $\tilde{\mathbb{D}}_i$  and they are not equivalent via the action of  $G_{36}$  if  $\mathcal{Y}_j$  and  $\mathcal{Y}_k$  do not belong to the same  $\tilde{\mathbb{D}}_i$ .

Let

$$\mathbb{D}_{(\alpha,\beta)_i} = \{\pi_\alpha^{-1}(\mathcal{Y}_\alpha) + (\pi_\beta^{-1}(\mathcal{Y}_\beta))^\omega \leq \mathbb{F}_2^{72} \mid \mathcal{Y}_\alpha, \mathcal{Y}_\beta \in \tilde{\mathbb{D}}_i, \omega \in C_{\text{Aut}(\mathcal{Y}_\beta(\chi))}(\langle \chi, \mu \rangle)\}.$$

**Remark 4.2.** Considering  $(\pi_\beta^{-1}(\mathcal{Y}_\beta))^\omega$  with  $\omega$  varying in  $C_{\text{Aut}(\mathcal{Y}_\beta(\chi))}(\langle \chi, \mu \rangle)$  is exactly the same as considering  $(\pi_\beta^{-1}(\mathcal{Y}_\beta))^\tau$  with  $\tau$  varying in a right transversal of

$$C_{\text{Aut}(\mathcal{Y}_\beta(\chi))}(\langle \chi, \mu \rangle)$$

in

$$C_{\text{Aut}(\mathcal{Y}_\beta(\chi))}(\langle \chi, \mu \rangle).$$

Obviously this makes the calculations faster.

**Lemma 4.3.** The code  $\mathcal{C}(\alpha) + \mathcal{C}(\beta) := \{v + w \mid v \in \mathcal{C}(\alpha) \text{ and } w \in \mathcal{C}(\beta)\}$  is equivalent, via the action of  $G$ , to an element of  $\bigcup_{i=1}^m \mathbb{D}_{(a,b)_i}$ .

*Proof.* By Lemma 4.1 and by construction of  $\bigcup_{i=1}^m \tilde{\mathbb{D}}_i$ , there exist  $i \in \{1, \dots, m\}$ ,  $\mathcal{Y}_a \in \tilde{\mathbb{D}}_i$  and  $\bar{\rho} \in G_{36}$  such that  $\pi_\alpha(\mathcal{C}(\alpha))^{\bar{\rho}} = \mathcal{Y}_a$ . Choose  $\rho \in \eta_\alpha^{-1}(\bar{\rho})$ . Then it is easy to observe that

- $\pi_\beta(\mathcal{C}^\rho(\beta))$  is a self-dual [36, 18, 8] code;
- $\langle \chi, \mu \rangle \leq \text{Aut}(\pi_\beta(\mathcal{C}^\rho(\beta)))$  (since  $\rho \in G$ );
- $(\pi_\beta(\mathcal{C}^\rho(\beta)))(\chi) = (\pi_\alpha(\mathcal{C}^\rho(\alpha)))(\chi) = \mathcal{E}_i$  (as in (1)).

Now,  $\{(\mathcal{Y}_\beta)^\tau \mid \mathcal{Y}_\beta \in \tilde{\mathbb{D}}_i, \tau \in C_{\text{Aut}(\mathcal{Y}_\beta(\chi))}(\langle \chi, \mu \rangle)\}$  is the set of all possible such codes, so  $(\pi_\beta(\mathcal{C}^\rho(\beta)))(\chi)$  is one of these codes.  $\square$

**Remark 4.4.** There are, up to equivalence in the full symmetric group  $S_{72}$ , only 22 codes in  $\bigcup_{i=1}^m \mathbb{D}_{(\alpha,\beta)_i}$  such that the minimum distance is at least 16, say  $\mathcal{D}_1, \dots, \mathcal{D}_{22}$ . They are all [72, 26, 16] codes. In particular

$$\dim(\mathcal{D}_i(\alpha) \cap \mathcal{D}_i(\beta)) = 10.$$

**Corollary 4.5.** The code  $\mathcal{C}(\alpha) + \mathcal{C}(\beta)$  is equivalent, via the action of the full symmetric group  $S_{72}$ , to a code  $\mathcal{D}_i$ , with  $i \in \{1, \dots, 22\}$ .

We can repeat in a completely analogous way all the procedure for the pairs  $(\alpha, \gamma)$  and  $(\beta, \gamma)$ , interchanging the roles of the elements  $\alpha, \beta$  and  $\gamma$ . Then we get the following.

**Corollary 4.6.** *The codes  $\mathcal{C}(\alpha) + \mathcal{C}(\gamma) := \{v + w \mid v \in \mathcal{C}(\alpha) \text{ and } w \in \mathcal{C}(\gamma)\}$  and  $\mathcal{C}(\beta) + \mathcal{C}(\gamma) := \{v + w \mid v \in \mathcal{C}(\beta) \text{ and } w \in \mathcal{C}(\gamma)\}$  are equivalent, via the action of the full symmetric group  $S_{72}$ , to some codes  $\mathcal{D}_j$  and  $\mathcal{D}_k$ , with  $j, k \in \{1, \dots, 22\}$ .*

This implies that

$$\dim(\mathcal{C}(\alpha) \cap \mathcal{C}(\gamma)) = 10 \quad \text{and} \quad \dim(\mathcal{C}(\beta) \cap \mathcal{C}(\gamma)) = 10. \quad (2)$$

Furthermore, by MAGMA calculations we get that

$$\dim(\mathcal{C}(\alpha) \cap \mathcal{C}(\beta) \cap \mathcal{C}(\gamma)) = 5. \quad (3)$$

Both statements can be verified by taking all the elements  $\alpha', \beta', \gamma'$  of order 2 and degree 72 in  $\text{Aut}(\mathcal{D}_i)$  such that  $\langle \alpha', \beta', \gamma' \rangle$  is conjugate to  $\langle \alpha, \beta, \gamma \rangle$  in  $S_{72}$ .

To get a contradiction it is now enough to observe that (2) and (3) are not compatible with the table in Lemma 3.3. So we conclude the following.

**Theorem 4.7.** *The automorphism group of a self-dual [72, 36, 16] code does not contain a subgroup isomorphic to  $C_2 \times C_2 \times C_2$ .*

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