# Metropolis-Hastings sampling of alternatives for route choice models 

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## Motivation

- a route choice model describes what way from an origin to a destination is chosen in a network
- universal set of alternatives is unknown and intractably large
- estimation of route choice models requires selection of a subset


## Approaches to route choice set generation

- modeling of consideration sets
- deterministic (e.g., K-SP) or stochastic (randomized SP)
- unrealistic: fail to capture the chosen alternative
- assume that decision maker considers all alternatives
- also unrealistic
- sampling protocol generates operational subset
- correct for sampling in the estimation


## Sampling of alternatives

- sample $\mathcal{C}_{n}$ with replacement from $\mathcal{C}$ according to $\{q(i)\}_{i \in \mathcal{C}}$
- add the chosen alternative
- $k_{i n}$ is the number of times alternative $i$ is contained in $\mathcal{C}_{n}$
- correct for sampling when estimating logit model

$$
P\left(i \mid \mathcal{C}_{n}\right)=\frac{e^{\mu V_{i n}+\ln \left(\frac{k_{i n}}{b(i)}\right)}}{\sum_{j \in \mathcal{C}_{n}} e^{\mu V_{j n}+\ln \left(\frac{k_{k n}}{b(j)}\right)}}
$$

where $\{b(i)\}_{i \in \mathcal{C}}$ is such that $q(i)=b(i) / \sum_{j \in \mathcal{C}} b(j)$
objective: sample paths according to pre-specified $\{b(i)\}_{i \in \mathcal{C}}$

## Using Markov chains (MCs)

- finite state space
- discrete time $k=0,1, \ldots$
- at time $k$, process is in state $i^{k}$
- $q(i, j)$ is one-step probability to reach state $j$ from state $i$
- process has a unique stationary distribution if
- every state eventually reaches every other state
- there is at least one state $i$ with $q(i, i)>0$
objective: build MC of routes with stationary distribution $\{q(i)\}_{i \in \mathcal{C}}$


## Metropolis-Hastings (MH) algorithm

- given a finite state space, positive weights $\{b(i)\}_{i}$ and "well-mixing" proposal transition distribution $q(i, j), \mathrm{MH}$ generates MC that converges to $q(i)=b(i) / \sum_{j} b(j)$

1. set iteration counter $k=0$
2. select arbitrary initial state $i^{k}$
3. repeat beyond stationarity
3.1 draw candidate state $j$ from $\left\{q\left(i^{k}, j\right)\right\}_{j}$
3.2 compute acceptance probability $\alpha\left(i^{k}, j\right)=\min \left(\frac{b(j) q\left(j, i^{k}\right)}{b\left(i^{k}\right) q\left(i^{k}, j\right)}, 1\right)$
3.3 with probability $\alpha\left(i^{k}, j\right)$, let $i^{k+1}=j$; else, let $i^{k+1}=i^{k}$
3.4 increase $k$ by one

## Application of MH for route choice set generation

- state space comprises $\mathcal{C}$
- weights $b(i)$ favor plausible paths (importance sampling)
- transition distribution $q(i, j)$ creates local path modifications
- too little variability: slow convergence
- too much variability: random search


## State space

- notation
- $\Gamma=$ a path (node sequence)
- Г $(u)=u$ th node of path 「
- $\Gamma(u, v)=$ sub-path from the $u$ th to the $v$ th node of $\Gamma$
- $|\Gamma|=$ number of nodes in path $\Gamma$
- state $=$ (a path $\Gamma$, two SPLICE locations, one SPLICE node)
- SPLICE locations $u, d \in \mathbb{N}$ with $1 \leq u<d \leq|\Gamma|$
- SPLICE node v
- (state expansion helps to compute $q(i, j)$ )


## Proposal transition distribution

- SPLICE operation
- compute (random) path from $\Gamma(u)$ to $v$
- compute (random) path from $v$ to $\Gamma(d)$
- replace $\Gamma(u, d)$ by that sequence
- SHUFFLE operation
- re-sample (uniformly) splice locations $u$ and $d$
- re-sample splice node $v$ near to to $\Gamma(u, d)$
- randomly select one procedure

$$
q(i, j)=\gamma q_{\text {SPLICE }}(i, j)+(1-\gamma) q_{S H U F F L E}(i, j)
$$

with $0<\gamma<1$

## Illustrative example

[old status quo]
destination
origin


## Illustrative example

[a SHUFFLE event occurs]
destination


## Illustrative example

[a SHUFFLE event occurs]
destination


## Illustrative example

[a SPLICE event occurs]
destination


## Illustrative example

[new status quo]
destination


## Discussion

- this eventually draws paths from any distribution
- computationally
- feasible path in every iteration
- need to reach and identify stationarity
- strong auto-correlation of subsequent paths
- behaviorally
- explorative travel behavior?
- occasional intermediate destinations?
- iterated DTA simulations (such as MATSim)
- an all-day plan is a generalized path
- small plan choice set for computational reasons
- challenge: capture distribution in simulated conditions


## Shortest path SPLICE reaches every path



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